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(54) Method and apparatus for data inversion in a memory device

(57) The present invention is a method of writing information to a synchronous memory device by examining a present word of N bits to be written, where each bit has a high or low value. The present word is compared to a previous word also having N bits to identify the number of bit transitions from a low value to a high value of vice versa. The present bit is inverted when the number of transitions is greater than N/2. To avoid the need for having an extra bit accompany data bytes to indicate the presence or absence of inversion, the present invention takes advantage of a data mask pin that is normally unused during writing operations to carry the inversion bit. Non-inverted data is written directly into the memory device while inverted data is first inverted again before writing to storage locations, so that true data is stored in the memory device.

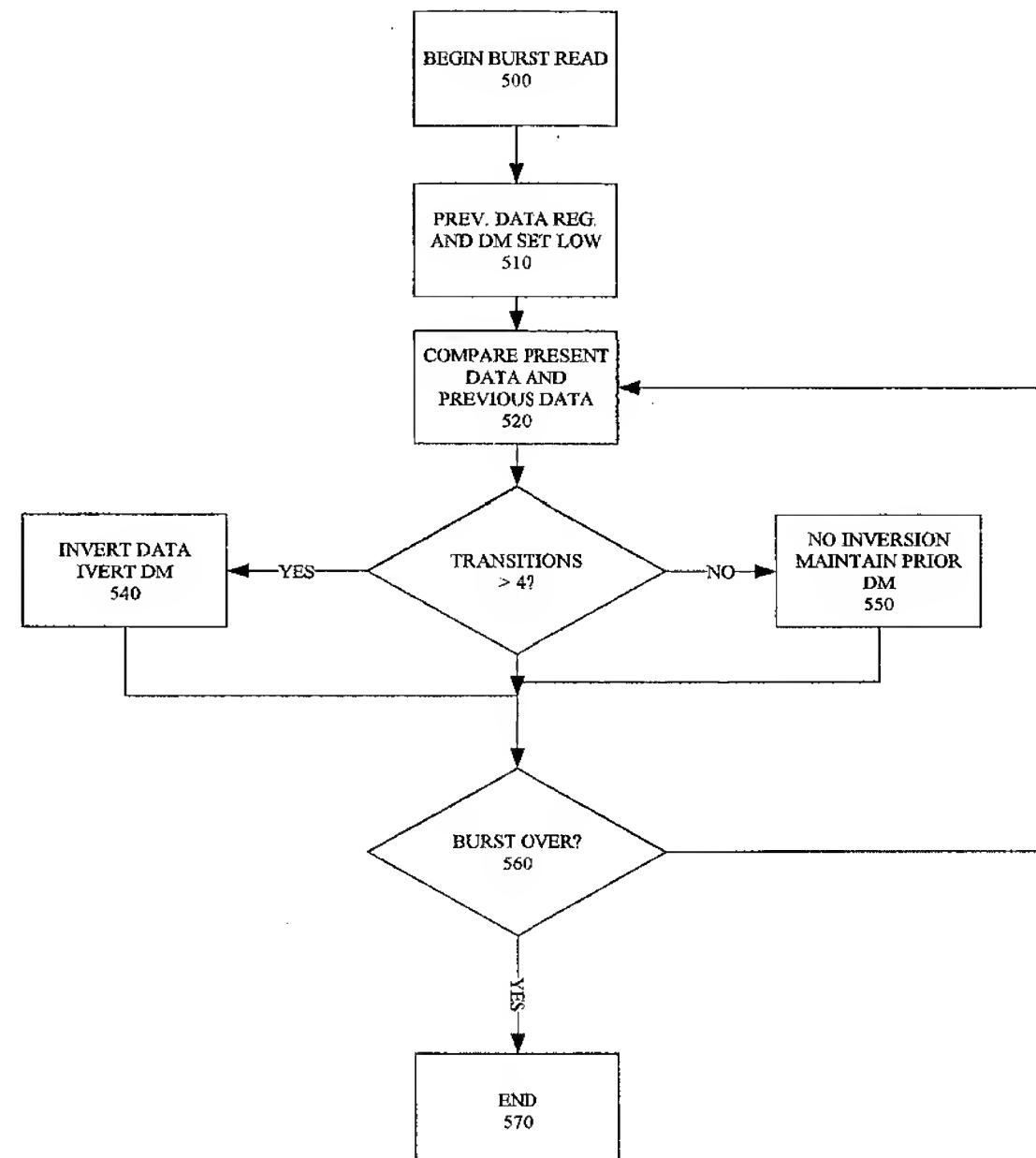


FIGURE 5

Description

[0001] The present invention relates to the field of data writing and reading.

[0002] Digital data is stored in memory devices as ones and zeros. Data is transferred into and out of a memory device via a conductive path called a "pin". Data is written into the memory as a string of high and low signals that represent the ones and zeros. Each time there is a transition between the low and high signals, the efficiency of the memory device is reduced. However, because data is somewhat random, there is no easy way to control the data to reduce the number of transitions. This problem can be understood by a review of writing to and from memory devices.

[0003] Figure 1 is a block diagram of a data source and a data receiver. The data source 100 communicates with the data receiver 102 via a bus 101. Data is transferred in blocks of 8 bits called "bytes". Each byte is sent to the data receiver 102 in parallel as 8 data signals on 8 lines with each signal line being high or low depending on the data being transmitted. For example, if a data value of one is sent, the bus line is at a high state. If the next signal sent is a zero, the bus line must be brought down to a low state to represent the new data. Then, if the succeeding signal is a one, the bus line must be again brought up to the high value. These changes of state are known as data transitions and require time and energy to achieve. Consider the following scenario, representing four bytes of data that are written to the data receiver 102.

Data Byte	# of Transitions
00000000	-
11111111	8
10001000	6
01110111	6

[0004] The total number of data transitions in the 24 bits of bytes 2, 3, and 4 in the above example is 20, or 83%. (Note that this assumes that each line of the bus was in a low state when the original byte is sent). This results in time and power penalties in the operation of the data transfer system.

30 Data Inversion Scheme

[0005] The prior art has provided a scheme to reduce the number of data transitions in writing to a memory device known as "data inversion". In the data inversion scheme, the number of data transitions between one byte and the next is determined and analyzed. If the number of transitions is greater than four, the data of the byte to be sent is inverted. In this way the number of transitions is made to be less than four. This scheme sets a bound of four on the greatest number of transitions that can occur from byte to byte, for a maximum of 50% transitions. The operation of the data inversion scheme is illustrated by reference to the table below.

Data Byte	Invert?	Data Byte Actually Sent	# of Transitions
00000000	-	00000000	-
11111111	Yes	00000000	0
10001000	No	10001000	2
01110111	Yes	10001000	0

[0006] After the first byte is sent, the next byte is examined to see how many transitions there would be if the byte was sent without inversion. Here, where the byte changes from all zeros to all ones, there would be 8 transitions. Since 8 is greater than four, the byte is inverted, becoming all zeros. When this byte is sent, there are no data transitions, since each line of the bus was already in the low state. Examining the next byte, it is compared to the actual byte sent (in this case the inverted byte) to determine the number of possible transitions. Without inversion, there will be two transitions, which is less than four, so the byte is sent without inversion. Comparing the next byte (01110111) to the byte that had been sent reveals 8 transitions if it is sent without inversion. Therefore the byte is inverted (to 10001000) and is sent with no transitions. Sending the same data using the inversion scheme results in only 2 transitions out of 24 bytes, or 8.3%. This improvement results in better and faster data transfer operation.

[0007] A disadvantage of the data inversion scheme is that it requires extra bits to be sent with the data byte so that the data can be properly used or stored at the data receiver, as well as extra bits when data is read from the receiver.

These extra bits require additional wiring for the bus, data sender and receiver, adding to the expense and complexity of a system.

[0008] The present invention is a method of writing information to a synchronous memory device by examining a present word of N bits to be written, where each bit has a high or low value. The present word is compared to a previous word also having N bits to identify the number of bit transitions from a low value to a high value of vice versa. The present bit is inverted when the number of transitions is greater than N/2. To avoid the need for having an extra bit accompany data bytes to indicate the presence or absence of inversion, the present invention takes advantage of a data mask pin that is normally unused during reading operations to carry the inversion bit. Non-inverted data is written directly into the memory device while inverted data is first inverted again before writing to storage locations, so that true data is stored in the memory device.

[0009] According to a first aspect of the present invention, a method of writing information to a synchronous memory device comprises:

examining a present word to be written, said present word having N bits, each of said bits having a low or high value;
 comparing said present word to a previous word having N bits and determining a number of transitions of bits from said low value in said previous word to said high value in said present word;
 inverting said present word when said number of transitions is greater than or equal to N/2;
 writing said present word to said synchronous memory device.

[0010] According to a second aspect of the present invention, a method of reading information from a synchronous memory device comprises:

examining a present word in said device to be read, said word having N bits, each of said bits having a high value or a low value;
 comparing said present word to a previous word having N bits and determining a number of transitions of bits having said low value in said previous word to bits having said high value in said present word;
 inverting said present word when said number of transitions is greater than or equal to N/2;
 reading said present word from said device.

[0011] The present invention will be described, by way of example, with reference to the accompanying drawings, in which:

Figure 1 is a block diagram of a data source and a data receiver;

Figure 2 is a diagram of an implementation for writing data to a device;

Figure 3 illustrates read circuitry for the DRAM of the present invention;

Figure 4 is an example of a DIM generator;

Figure 5 is a flow diagram that illustrates the operation of the invention during a burst read operation;

Figure 6 is a block diagram of the invention used with a microprocessor interface;

Figure 7A is a block diagram of an alternate embodiment of a read block of the present invention; and,

Figure 7B is an alternate embodiment of a write block of the present invention.

[0012] The present invention is a method and apparatus for reading and writing data. In the following description, numerous specific details are set forth to provide a more thorough description of embodiments of the invention. It is apparent, however, to one skilled in the art, that the invention may be practiced without these specific details. In other instances, well known features have not been described in detail so as not to obscure the invention.

[0013] Currently there has not been provided an efficient scheme for using data inversion in a dynamic random access memory (DRAM) system. One of the advantage of using DRAMs in a system is the low cost of DRAMs and the commodity nature of such memory devices. A DRAM communicates with other devices via "pins" that extend from the DRAM package. If a DRAM using a data inversion scheme requires extra pins for the inversion information, the DRAM might not be usable in many applications where the geography is for DRAMs having fewer pins. Similarly, the expense of building and selling a DRAM with an extra pin could make it commercially unviable.

[0014] The present invention takes advantage of an existing pin called the data mask (DM) pin that is associated with the data pins of a DRAM. The DM pin is typically used on data input but not on data output. The invention takes advantage of the DM pin on data output to provide a signal indicating the presence or absence of data inversion when the data is being outputted by the DRAM.

5 [0015] In the DRAM of the preferred embodiment of the invention, there are 32 data bits DQ0 to DQ31 divided into four bytes each DQ0 - DQ7, DQ8 - DQ15, DQ 16 - DQ23, and DQ 24 - DQ31. Each of the data bytes has an associated data mask pin DM0 - DM3 that is used during write operations to mask the data input as appropriate. When any of DM1 to DM3 are high, the data input at the same timing are masked.

10 Write Operations

[0016] Write operations to the DRAM must include an indication of whether the incoming data is inverted or not inverted. A data inversion mask (DIM) bit is included with the data to be written so that the DRAM can store the data in its true state. One implementation for accomplishing this is illustrated in Figure 2. Incoming data is provided as input 201 to an exclusive OR gate 203, along with DIM bit 202. When the DIM bit is high, the data coming in has been inverted so it is re-inverted via XOR gate 203 before being stored in the data array 204.

Read Operations

20 [0017] Figure 3 illustrates read circuitry for the DRAM of the present invention. Each byte of data has an associated DM pin. For example, data byte D[7:0] is associated with DM0, byte D[15:8] is associated with DM1, byte [23:16] is associated with DM2, and byte D[31:24] is associated with DM3. During read operations, the data is provided from array 300 directly to the data output pins and to a DIM generator such as generators 301.1 through 301.4. The DIM generator determines of the data should be inverted or not to reduce data transitions during the read operation. The 25 output of the DIM generator is outputted on the DM pin associated with the data byte.

[0018] An example of a DIM generator is illustrated in Figure 4. A last data register 401 stores the most recently sent byte (as it was actually sent, inverted or non inverted). A new data register 402 stores the current data byte to be sent. The outputs of registers 401 and 402 are provided to comparator 403. If the number of transitions is greater than 4 bits, the comparator 403 outputs a high signal to be output on the DM pin with the data byte.

30 Burst Operation

[0019] Often reads are done from a DRAM in a burst mode. Figure 5 is a flow diagram that illustrates the operation of the invention during a burst read operation. At block 500 the burst read begins. At block 510 the previous data 35 register and the DM signal are set to low. At block 520 the present data and previous data are compared to determine the number of transitions that will occur when sending the present data. At decision block 530 the number of transitions is checked to see if it is greater than 4. If yes, the data is inverted and the DM is inverted at block 540. If the result at decision block 530 is no, the data is not inverted and the previous DM value is maintained at block 550.

[0020] After either of block 540 or 550, a determination is made at decision block 560 as to whether the burst operation 40 has ended. If yes, the process ends at block 570. If no, the present data is stored in the prior data register, the new data byte is obtained, and the process returns to block 520.

Pin Advantages

45 [0021] The present invention has the advantage of providing the ability to provide data inversion for a DRAM without the need for excessive increased pin count, especially for read operations. By using the a pin that is normally unused during read operations, i.e. the DM pin, the invention permits byte reads using data inversion with no additional pins for the read operation.

50 Termination

[0022] As described above with respect to the burst read example, an assumption was made that the bus was in a low state when the burst read operation began. This is not required and the invention is not limited to such a scheme or assumption. The termination of the bus can be any one of a number of ways as described below.

55

A. Pulled Low -

[0023] In one embodiment of the invention, the bus is gradually pulled low after reads. One method of accomplishing

this is described in a co-pending patent application entitled NMOS Termination, assigned to the assignee of the present invention, and filed on _____, serial number _____. In this scheme, NMOS transistors are coupled to the bus lines and gradually pull them to ground over some period of time. When this scheme is used, a read operation is assumed to begin with all lines low and the number of transitions for the first byte of data is determined based on this assumption.

B. Pulled High -

[0024] In another embodiment, the bus can be pulled high by a PMOS termination scheme so that over time, after activity on the bus, all lines are pulled to a high state. The assumption in this scheme is that all lines are high at the beginning of a read operation and therefore the number of transitions for the first byte is determined with respect to an all high state.

C. Midpoint Termination -

[0025] In an alternative embodiment, the bus lines are pulled to a midlevel condition over time. When this scheme is used, the first byte in a read operation is sent non-inverted since there will be no difference in pulling the bus lines high or low from the midlevel state.

D. Unterminated Unlatched -

[0026] In this embodiment, the bus lines are not terminated at all and a first data byte is always sent unterminated. This results in lost efficiency for the first byte, since it may be possible to have greater than four transitions for the first byte.

E. Unterminated Latched -

[0027] This embodiment leaves the bus lines in the same state as the most recent byte on the bus and latches that value. For a subsequent operation, the previous byte register contains the latched value and the first byte is compared to that latched value to determine the number of transitions.

Processor Interface Logic Application

[0028] Although described in connection with a DRAM, the present invention has other applications as well. For instance, the scheme can be used in connection with a microprocessor interface logic to allow a data inversion scheme to be used. An example of such an embodiment is illustrated in Figure 6.

[0029] A DRAM 601 includes, for example first and second data banks 601A and 601B each having an output bus 602A and 602B along with associated DIM signals 603A and 603B (note that the DIM signals could be provided on DM pins). The outputs of the DRAM are provided to memory controller 604. Within the memory controller 604 the data outputs 602A and 602B of the DRAM are combined in a 16 bit bus 605 to microprocessor 608. The DIM signals 603A and 603B are coupled as inputs to an XOR gate 606 to produce an output signal 607 to microprocessor 608.

Power Optimization

[0030] An embodiment of a read block of the present invention is illustrated in Figure 7A. This embodiment provides DC power optimization when the I/O is open drain (i.e. the bus is terminated to a logical 1). Referring to Figure 7A, read data 701 (8 bits) is coupled to flip-flop 702, comparator 703, and NOR gate 704. A load signal 705 is coupled to the enable input of flip-flop 702. The output of flip-flop 702 is coupled to one input of OR gate 706 along with DC optimizing signal 707. The DC optimizing signal is a logical 1 which is OR'ed into the previous data field of the comparator block 703. This OR'ing guarantees that the output of the comparator will have more than half of its values being a logical 1, saving DC power. The comparator 703 compares the present data field with the previous data field to see if there are more than 4 bits with different values. If so, the output (DIM) is a logical 1 and is provided to NOR gate 704 and appropriate inversion of the data is provided.

[0031] If the bus is terminated to a logical 0, then the DC optimizing signal can be a logical 0 as well, so that more than half of the outputs are biased to 0, also saving power.

[0032] Referring now to Figure 7B, a block diagram of a write block is shown. Write data 710 (e.g. 32 bits) is provided to flip-flop 711 and a DIM signal 712 is provided to flip-flop 713. Both flip-flops are clocked with write clock signal 714. The outputs are coupled as inputs to NOR gate 715 to produce the data to be written into the memory array.

[0033] The present invention has been described in terms of the transmission of data but has equal application to any digital signals including address or command signals.

[0034] Thus, a method and apparatus for data inversion in memory devices has been described.

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Claims

1. A method of writing information to a synchronous memory device comprising:

10 examining a present word to be written, the present word having N bits, each of the bits having a low or high value;
 comparing the present word to a previous word having N bits and determining a number of transitions of bits from the low value in the previous word to the high value in the present word, or from the high value in the previous word to the low value in the present word;
 15 inverting the present word when the number of transitions is greater than or equal to N/2;
 writing the present word to the synchronous memory device.

2. A method of reading information from a synchronous memory device comprising:

20 examining a present word in the device to be read, the word having N bits, each of the bits having a high value or a low value;
 comparing the present word to a previous word having N bits and determining a number of transitions of bits having the low value in the previous word to bits having the high value in the present word, or the high value in the previous word to bits having the low value in the present word;
 25 inverting the present word when the number of transitions is greater than or equal to N/2;
 reading the present word from the device.

3. The method of claim 1 or claim 2, further including sending a data inversion mask (DIM) bit with the present word to indicate when the present word has been inverted.

30 4. The method of claim 3, wherein the DIM bit is sent on an existing pin on the memory device.

5. The method of claim 4, wherein the existing pin is a data mask (DM) pin.

35 6. The method of any one of the preceding claims, wherein the previous word is presumed to have all low values when the previous word is unknown.

7. The method of any one of claims 1 to 5, wherein the previous word is presumed to have all high values when the previous word is unknown.

40 8. The method of any one of claims 1 to 5, wherein the present word is not inverted when the previous word is unknown.

9. The method of any one of the preceding claims, wherein the present word is latched and used as the previous word in a subsequent read operation.

45 10. The method of any one of the preceding claims, wherein the information comprises data, such as graphical data.

11. The method of any one of the preceding claims, wherein the information comprises one or more commands, or address information.

50 12. The method of any one of the preceding claims, wherein the synchronous memory device comprises a DRAM.

13. The method of any one of the preceding claims, wherein the system includes biasing means to bias the previous word so that the present word is biased to have more of one logical value whether or not the present word is inverted.

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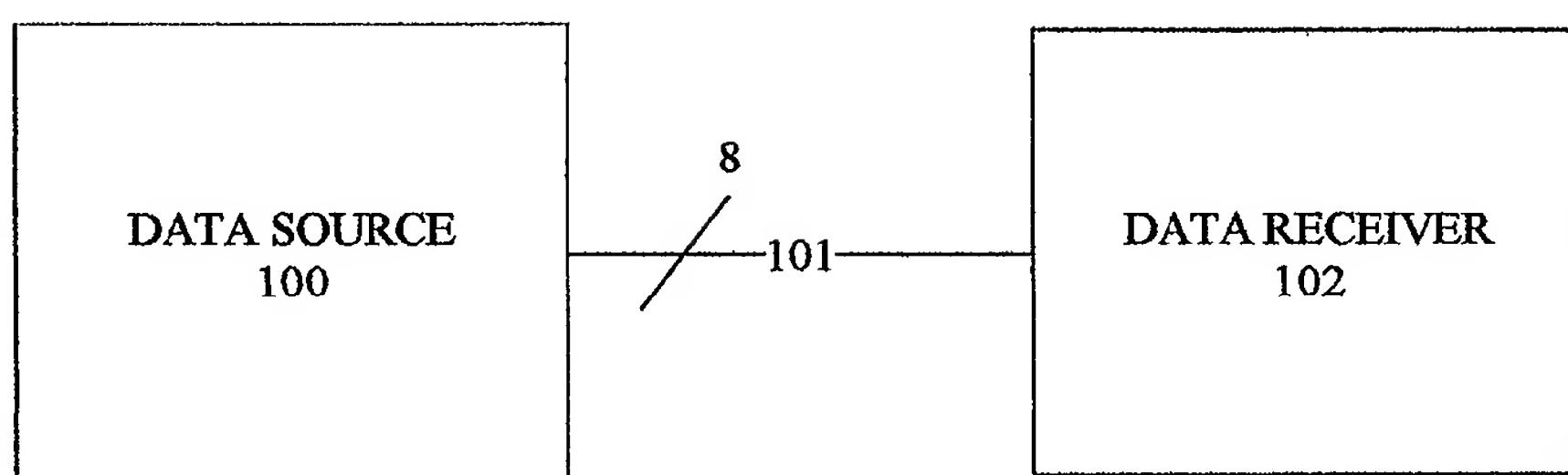


FIGURE 1

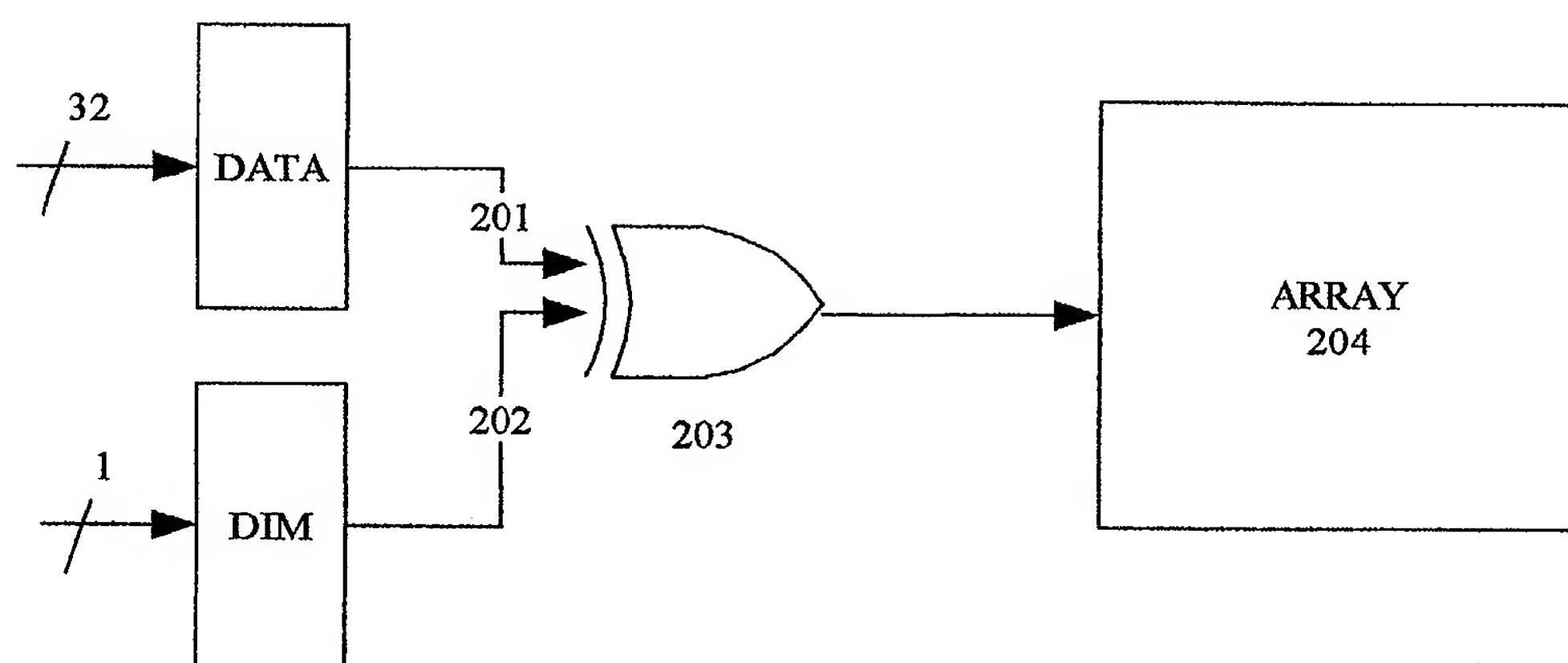


FIGURE 2

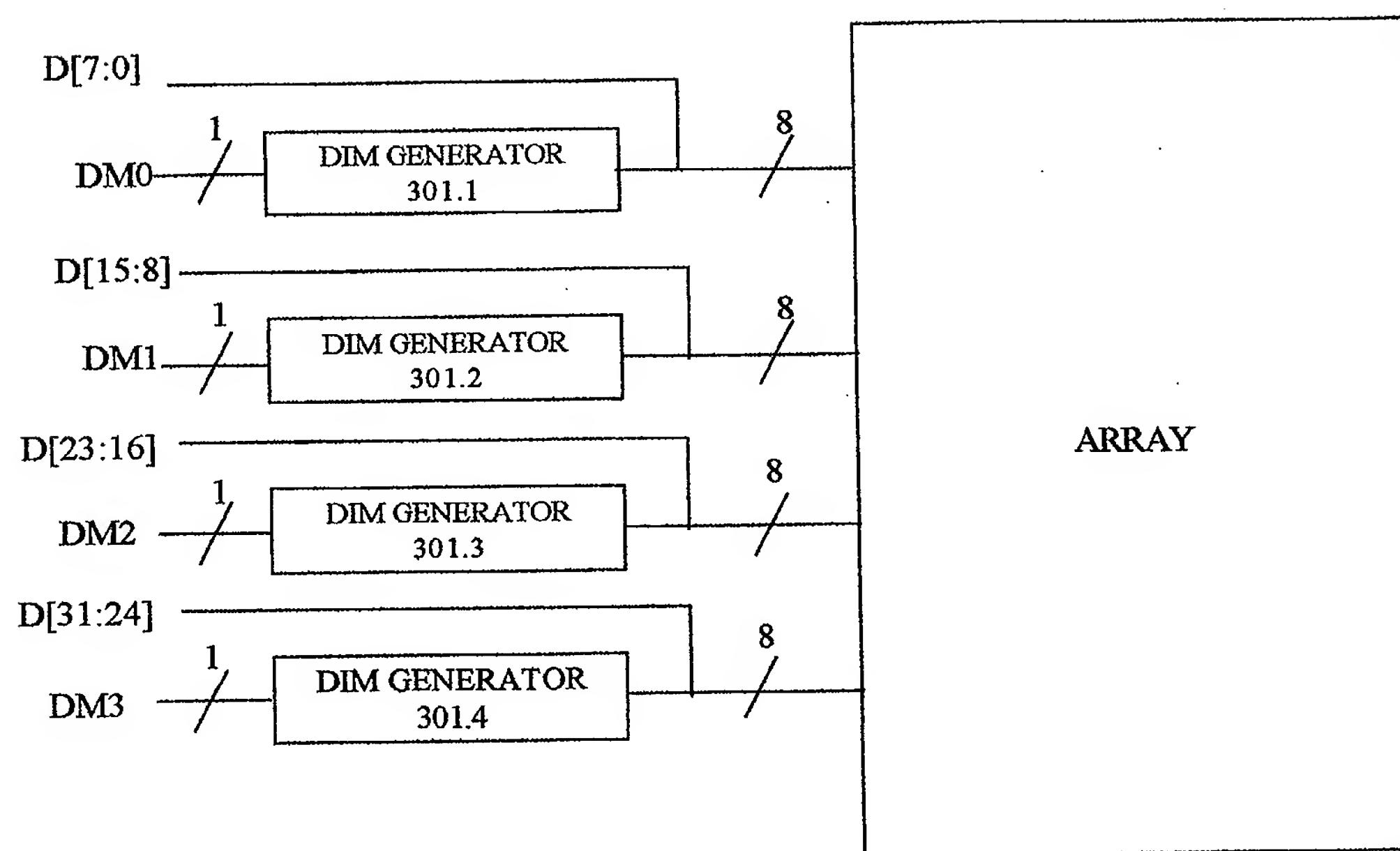


FIGURE 3

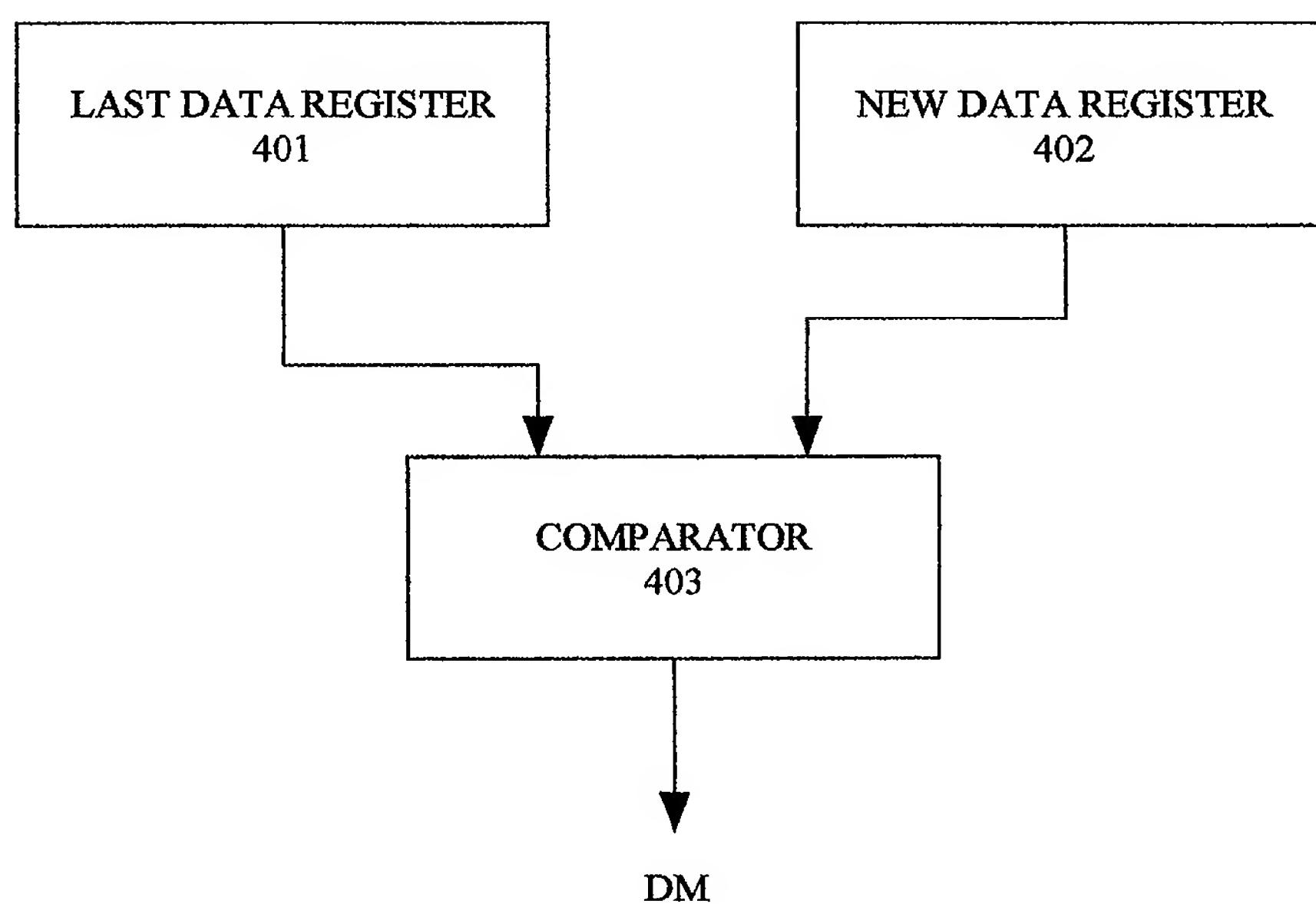


FIGURE 4

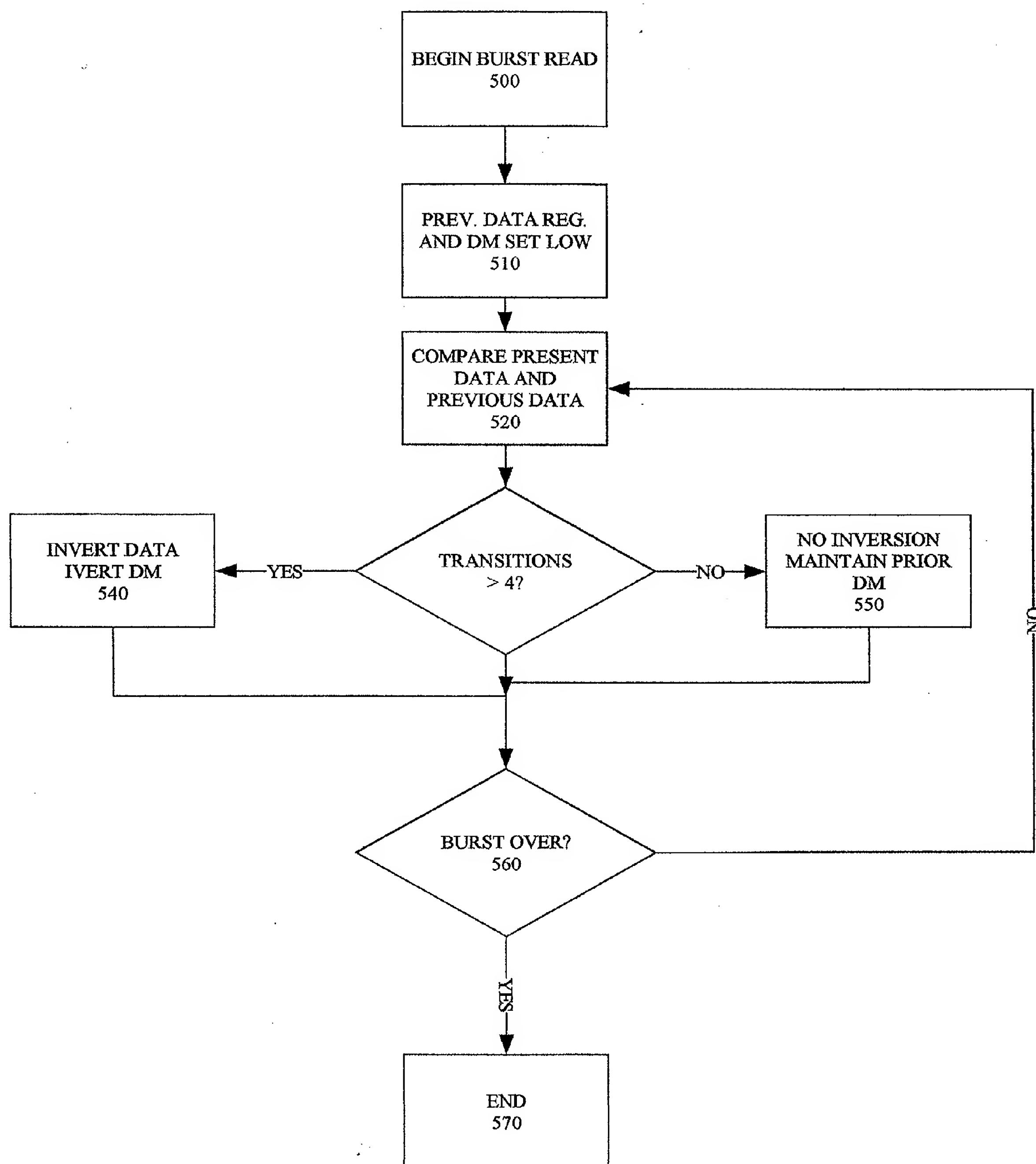


FIGURE 5

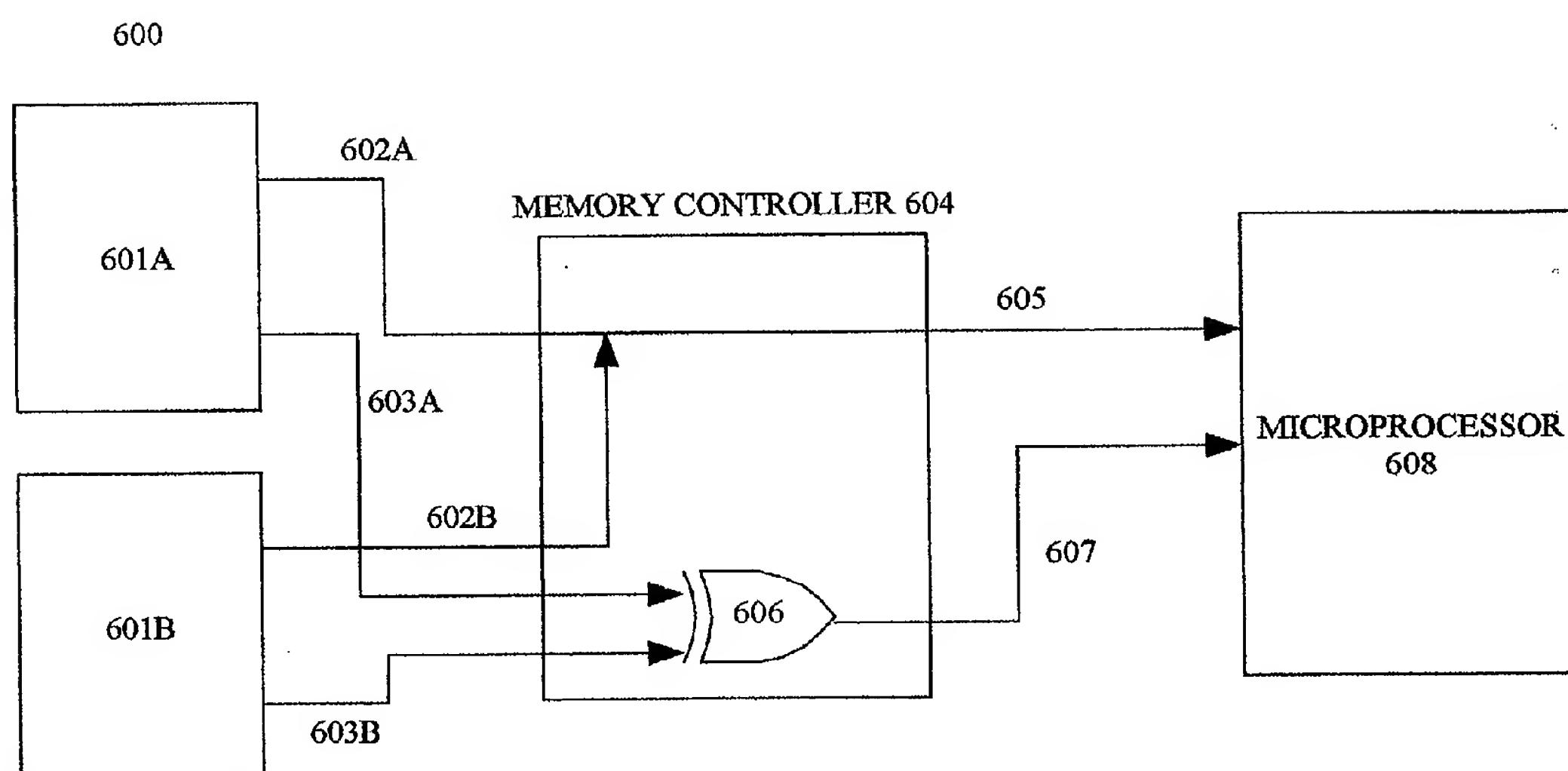
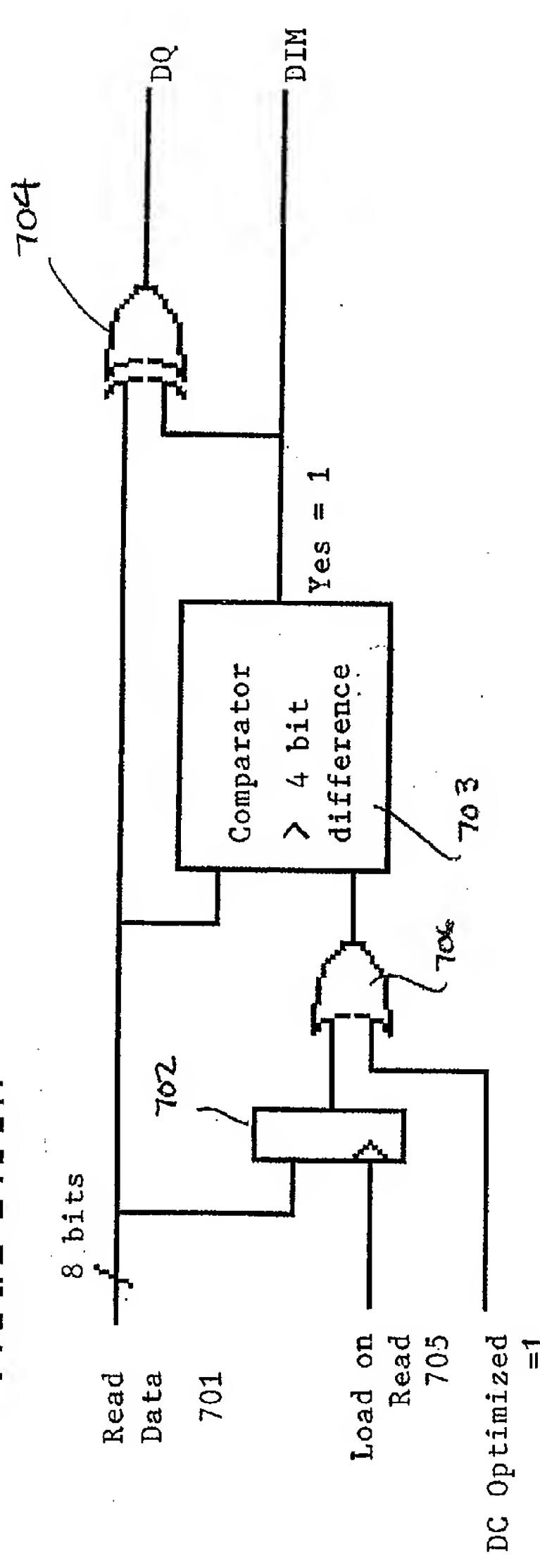


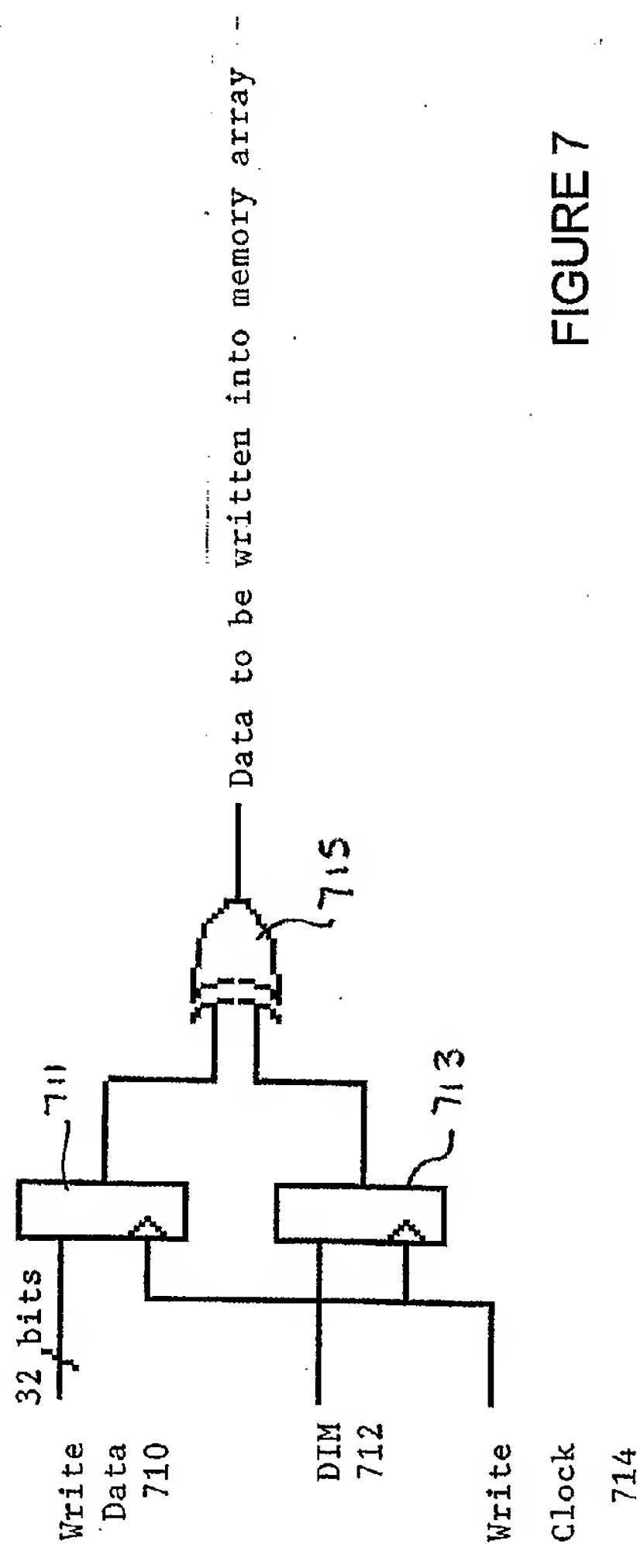
FIGURE 6

Simple block diagram of data integration

Read block



■ Write block





EUROPEAN SEARCH REPORT

Application Number
EP 03 25 0820

DOCUMENTS CONSIDERED TO BE RELEVANT			
Category	Citation of document with indication, where appropriate, of relevant passages	Relevant to claim	CLASSIFICATION OF THE APPLICATION (Int.Cl.7)
X	US 6 046 943 A (WALKER DARRYL G) 4 April 2000 (2000-04-04) * abstract * * column 1, line 14 - line 19 * * column 1, line 66 - column 2, line 21 * ---	1-13	G11C7/10
X	US 5 748 902 A (DALTON SCOTT W ET AL) 5 May 1998 (1998-05-05) * abstract * * column 1, line 65 - column 2, line 36 * * column 3, line 19 - line 40 * * column 5, line 44 - column 7, line 6; figures 2,4 *	1,2	
X	US 4 667 337 A (FLETCHER RODERICK J) 19 May 1987 (1987-05-19) * abstract *	1,2	
X	US 6 335 718 B1 (HONG JIN CHEOL ET AL) 1 January 2002 (2002-01-01) * abstract * * column 5, line 36 - column 6, line 32 * * column 6, line 65 - column 7, line 59; figures 3-12 *	1,2	<p>TECHNICAL FIELDS SEARCHED (Int.Cl.7)</p> <p>G11C</p>
The present search report has been drawn up for all claims			
Place of search	Date of completion of the search	Examiner	
MUNICH	30 April 2003	Ramcke, T	
CATEGORY OF CITED DOCUMENTS		T : theory or principle underlying the invention E : earlier patent document, but published on, or after the filing date D : document cited in the application L : document cited for other reasons & : member of the same patent family, corresponding document	
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ON EUROPEAN PATENT APPLICATION NO.

EP 03 25 0820

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30-04-2003

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